



# Comparing Inclusion Techniques on Chemical Engineering Problems

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# Outline

- Affine arithmetic
- Root-finding procedure
- Numerical examples
  - affine vs. interval arithmetic
  - convex envelopes
- Conclusions



# Affine arithmetic

# Dependency problem

- The simple example

$$X - X \neq [0, 0] \quad (\text{diameter: double of } X)$$

- Affine arithmetic keeps track of first-order correlation between computed and input quantities
- For example, with affine arithmetic (Figueiredo'97)

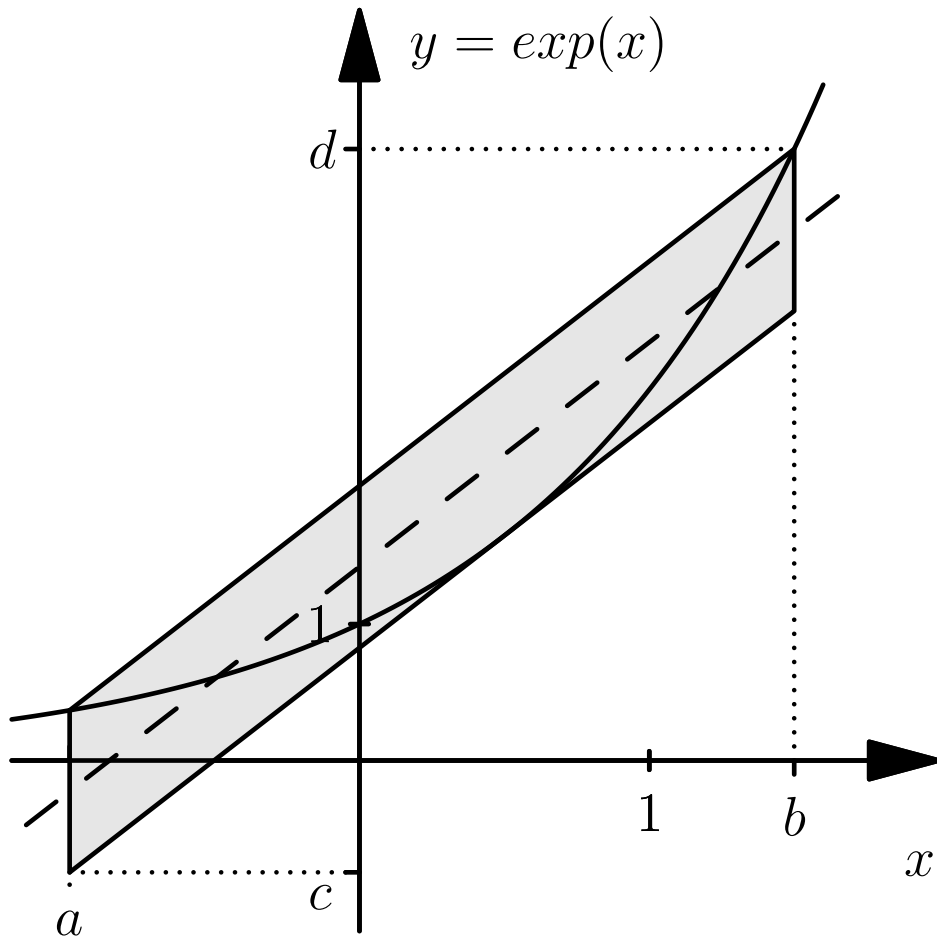
$$X = x_0 + x_1 \cdot \varepsilon_1 \quad -1 \leq \varepsilon_1 \leq 1$$

$$(x_0 + x_1 \cdot \varepsilon_1) - (x_0 + x_1 \cdot \varepsilon_1) = [0, 0]$$

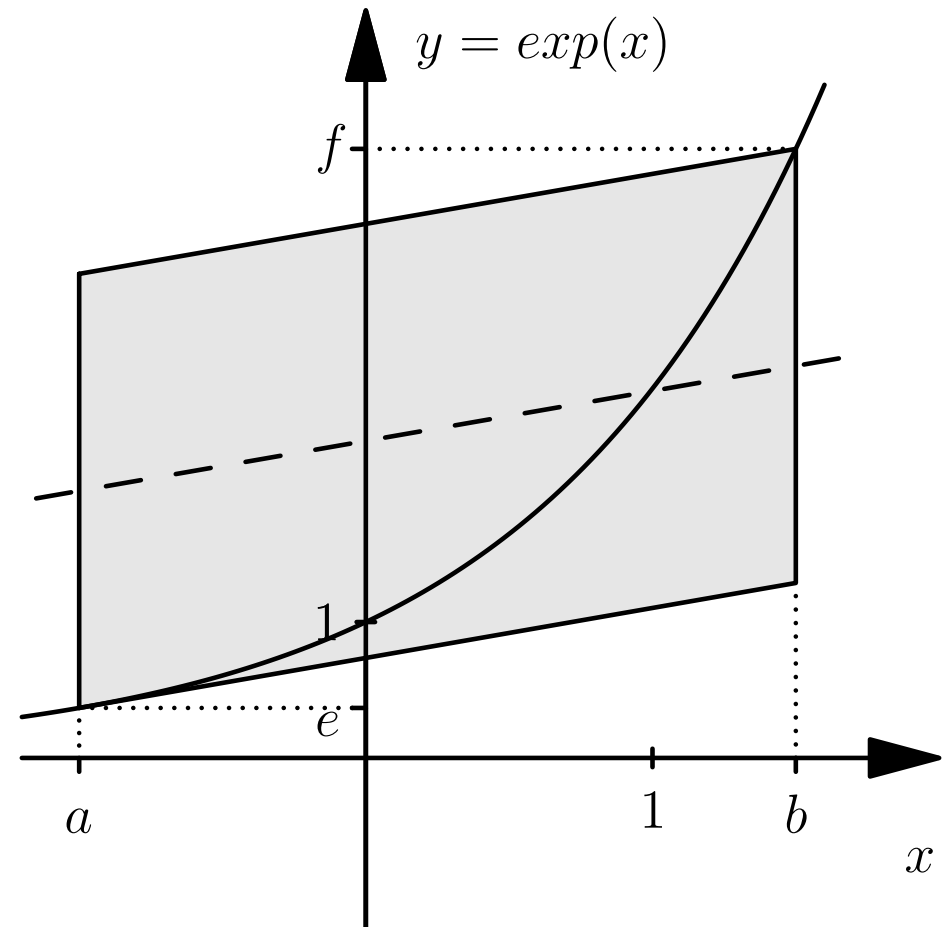
- Similarly, affine operations are ideally treated (apart from the rounding error)

# Non-affine operations

A new (unused) variable is introduced to rigorously enclose the approximation error



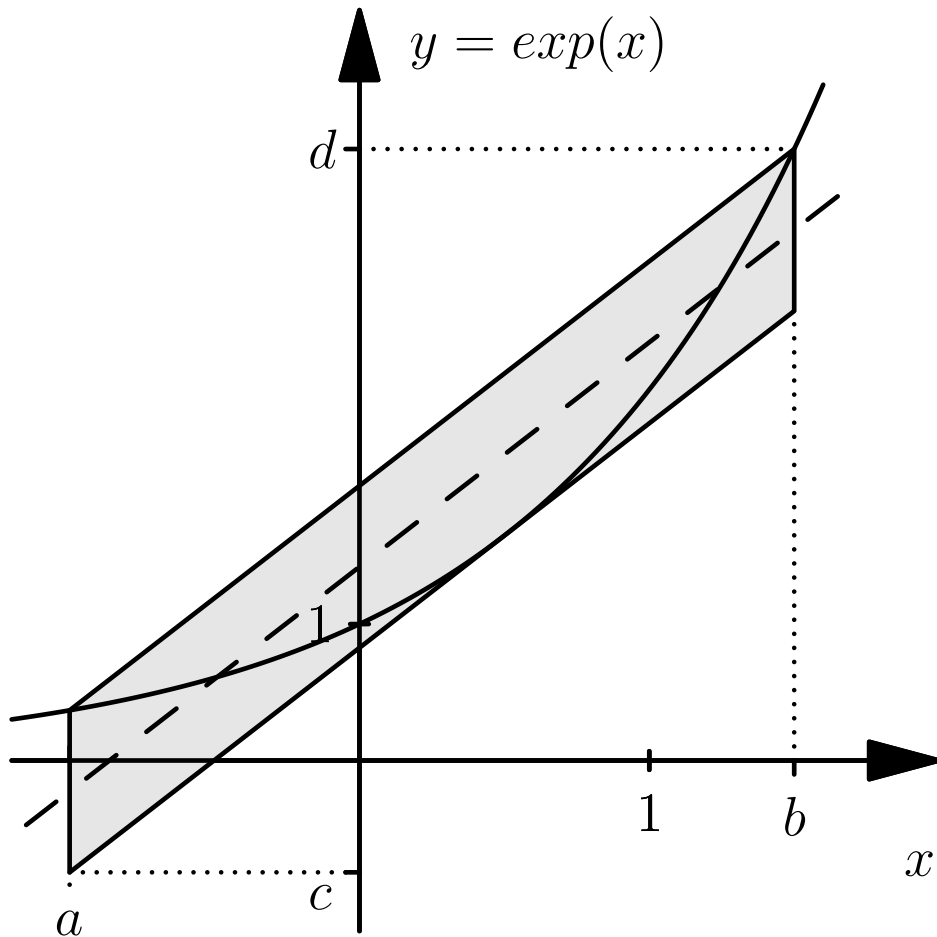
minimax



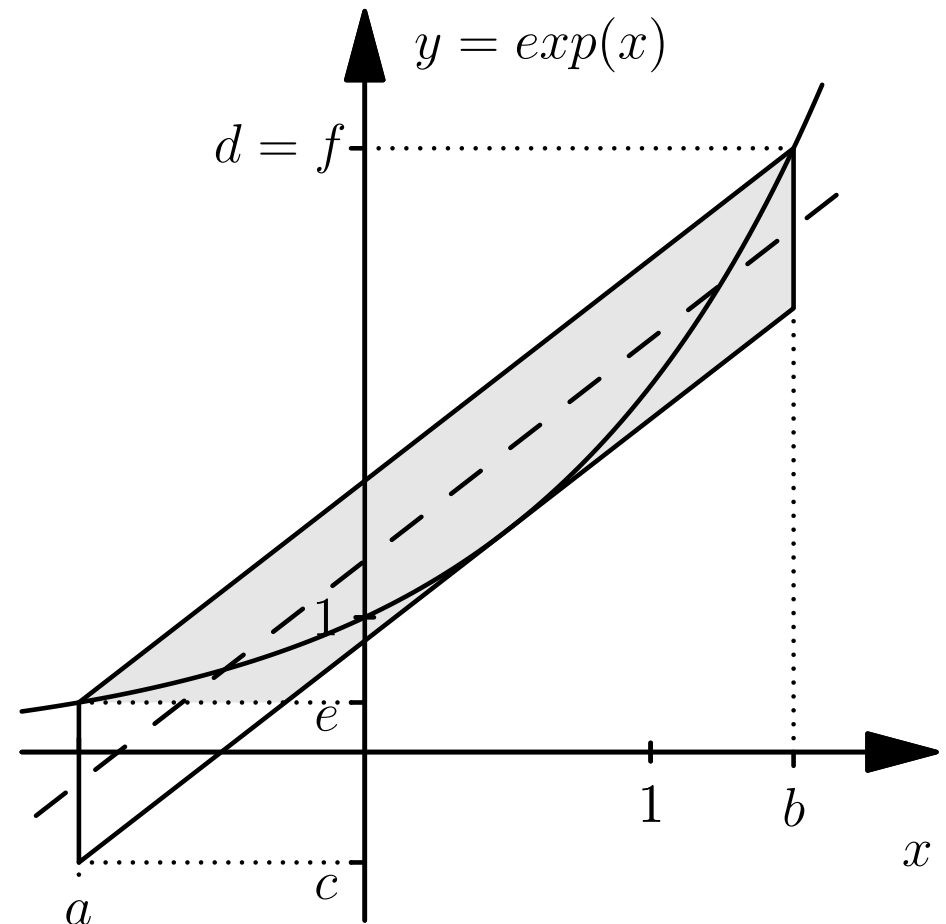
minrange

# Non-affine operations (continued)

A new (unused) variable is introduced to rigorously enclose the approximation error



minimax



mixed AA/IA

# Binary operations

- With ordinary **interval arithmetic**

$A(B+C) \in AB+AC$  (subdistributivity)

$X/X \neq [1, 1]$

- The modified **affine arithmetic** of Kolev'04

$A(B+C) = AB+AC$  if  $B$  and  $C$  are independent

$X/X = [1, 1]$

- Optimal multiplication w.r.t. range Kolev'07
- Optimal multiplication w.r.t. width Miyajima'03



# Root-finding procedure



# Major components

- **Linearization**

ILA\* or LIA\*

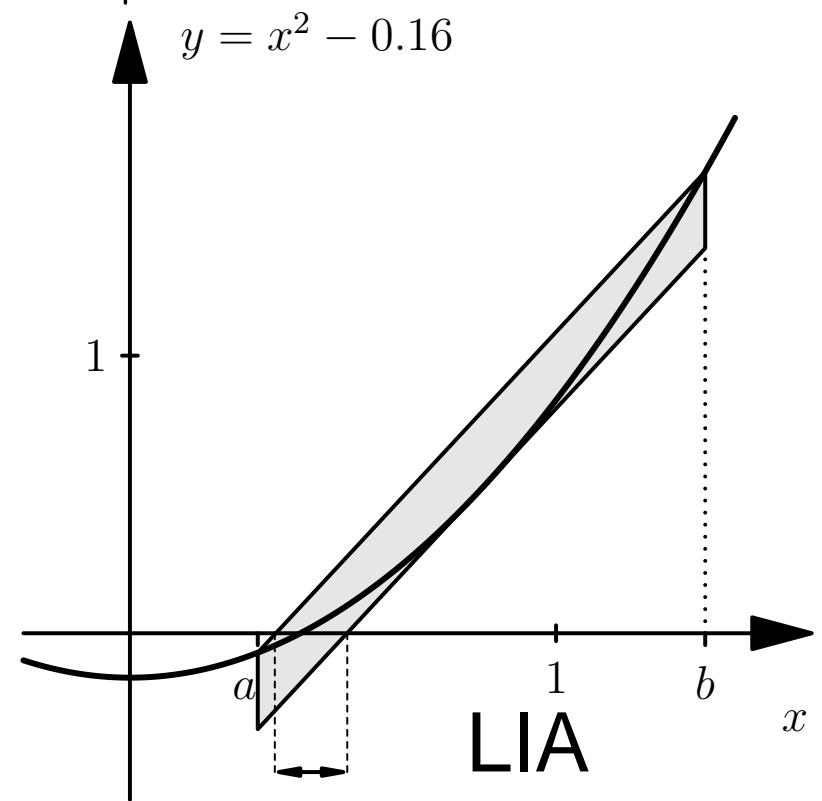
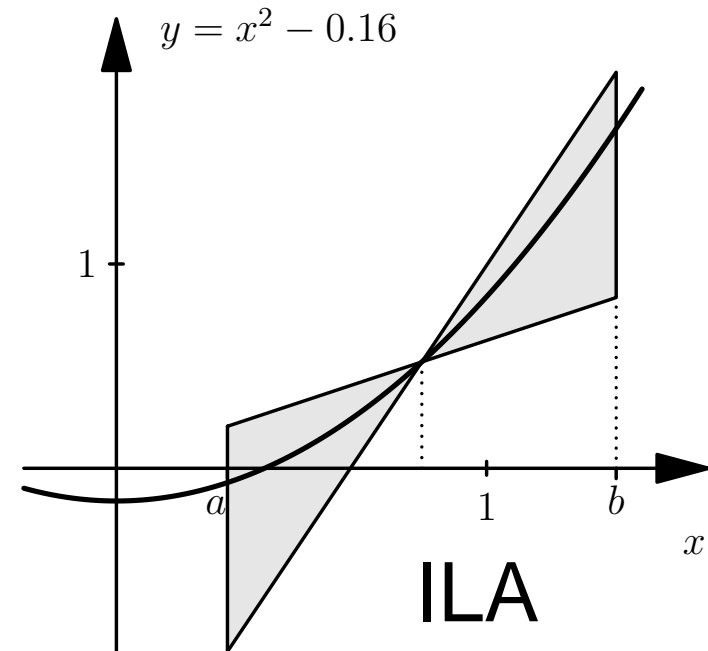
- **Pruning the box**

based on CP, LP

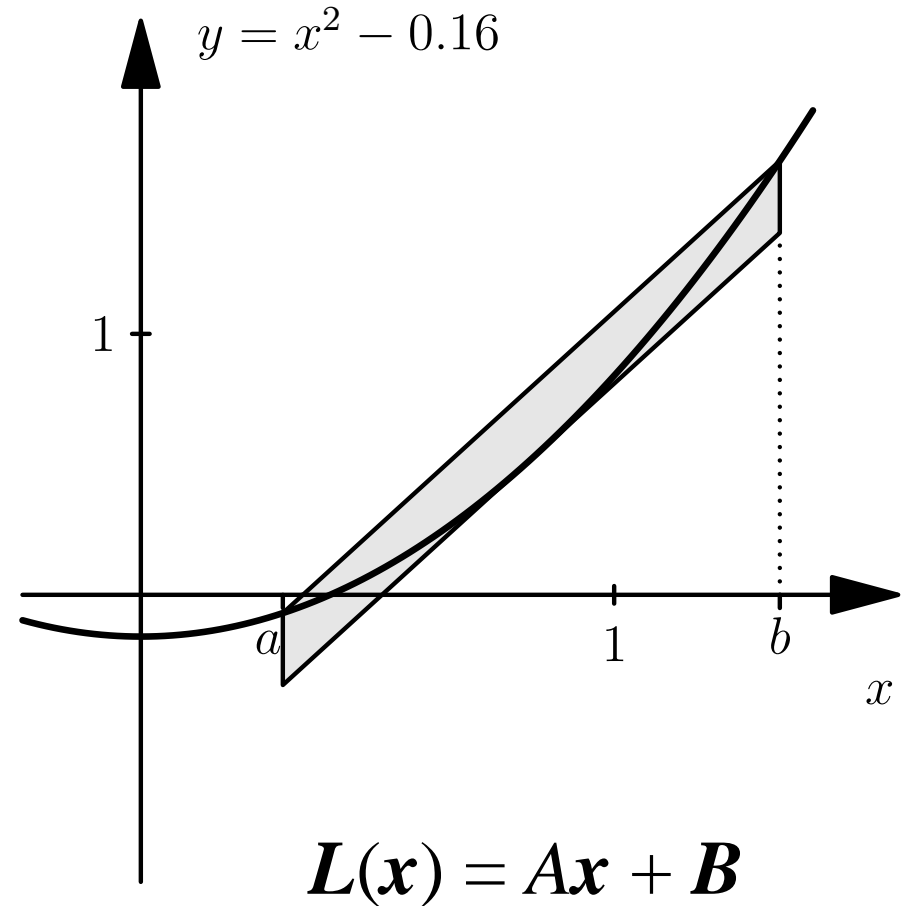
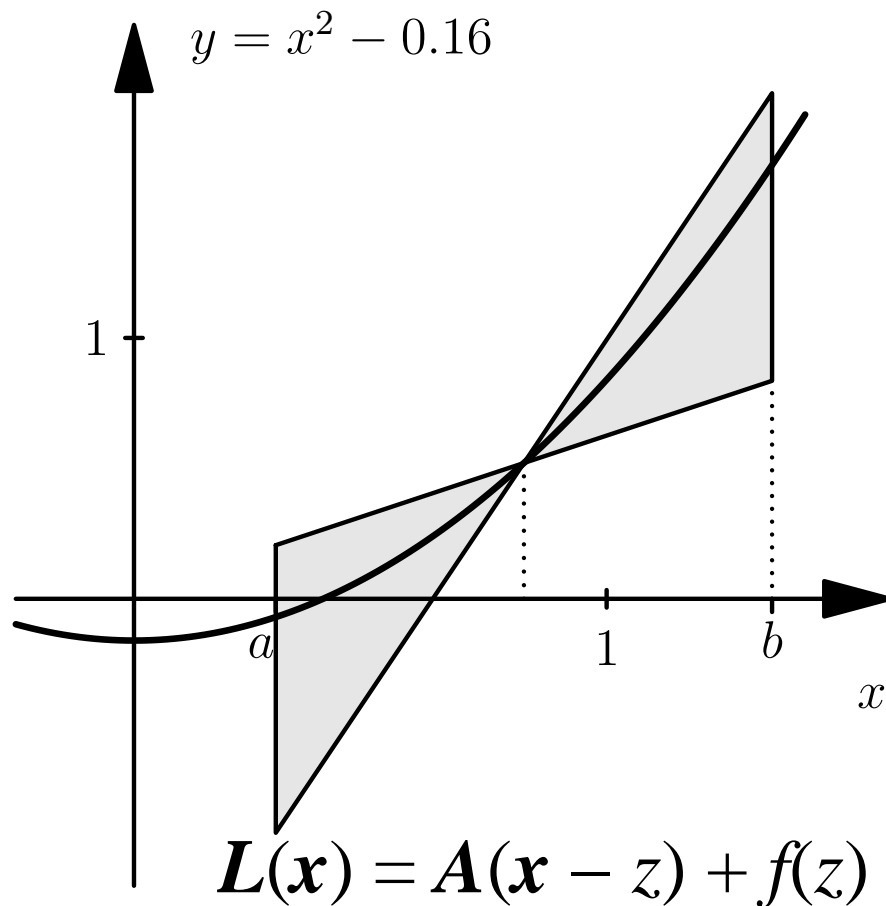
- **Splitting the box**

(not discussed here)

\* Kolev'06



# Interval Linear Approximation    Linear Interval Approximation



## Advantages of LIA

- Keeps track of dependencies (coefficient matrix)
- The solution set has a simpler form (convex)
- LP is directly applicable

# Pruning

$$L(x) = Ax + B \quad x \in X$$

- Hull solution is straightforward

$$Y = -A^{-1}B$$

$$X^{new} = X \cap Y$$

- In the LP pruning, the following  $2n$  LP subproblems have to be solved

$$\min/\max \quad x_j \quad \text{for all } j$$

subject to

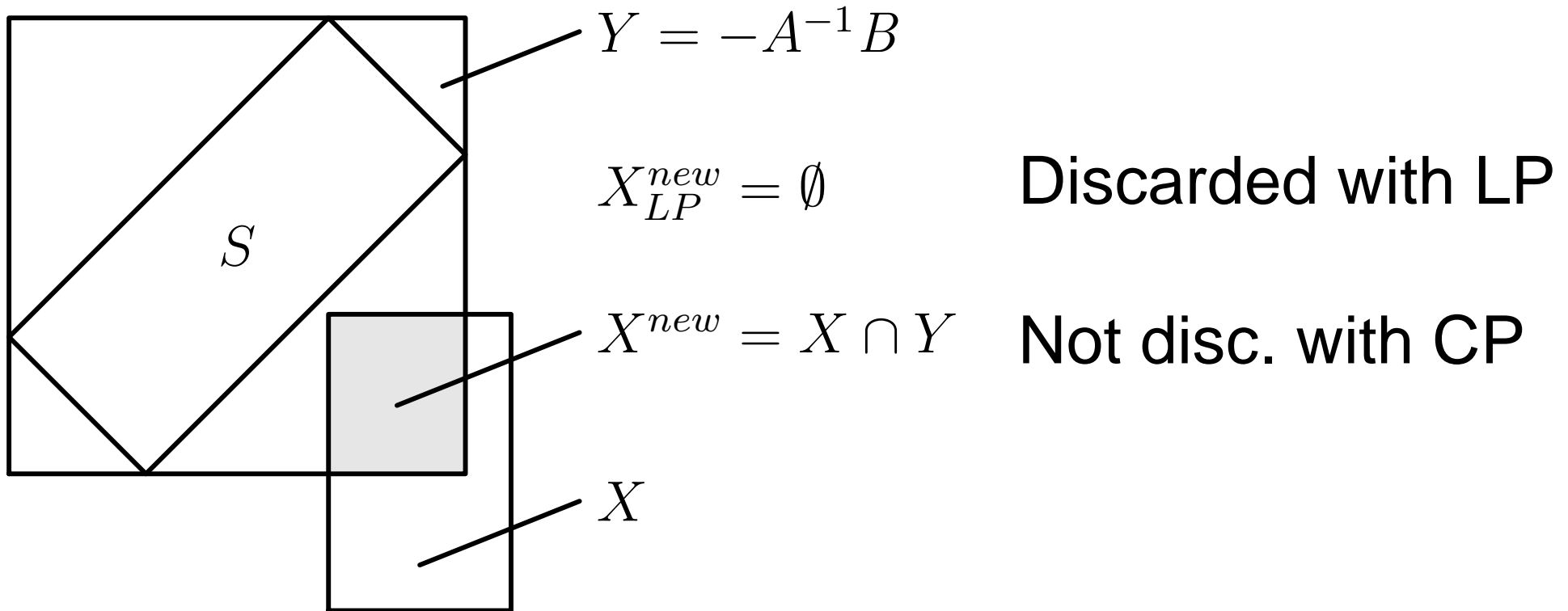
$$-b_U \leq Ax \leq -b_L$$

$$x_L \leq x \leq x_U$$

# Comparing the pruning techniques I

$$L(x) = Ax + B \quad x \in X$$

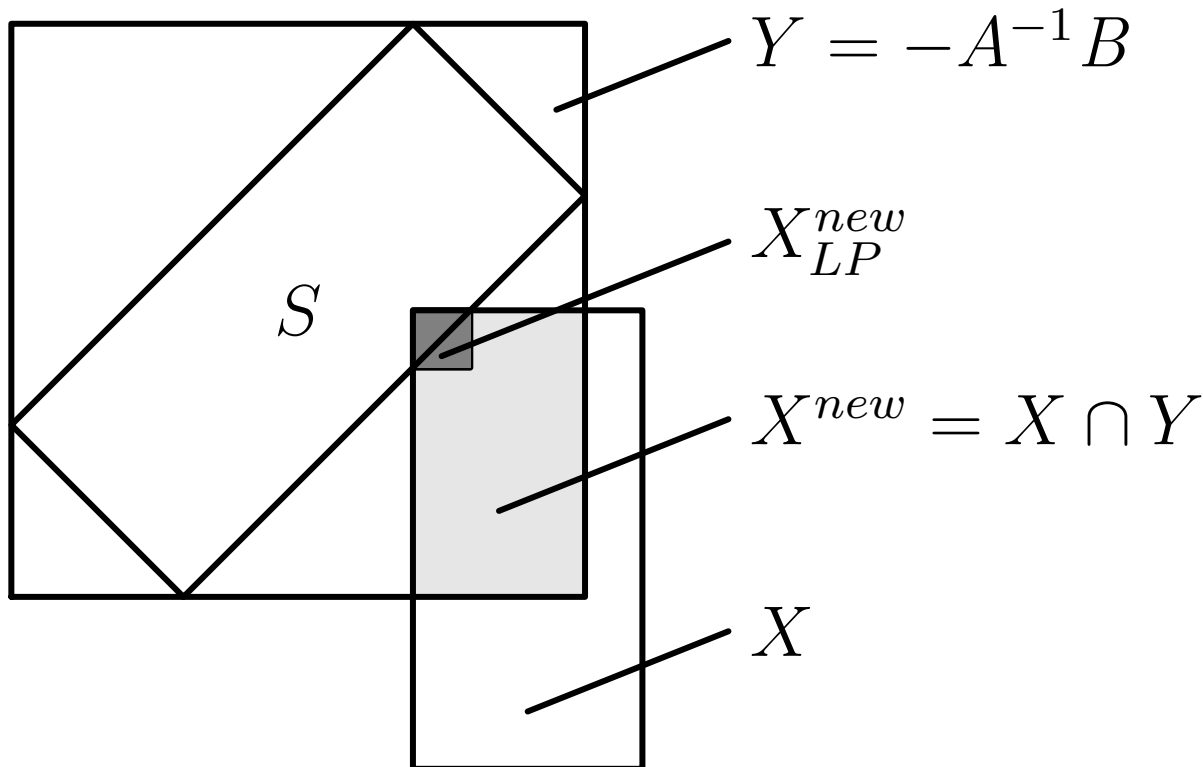
No solution in  $X$



# Comparing the pruning techniques II

$$L(x) = Ax + B \quad x \in X$$

There may be a solution in  $X$



Much smaller part is kept with LP pruning

# Revised LP pruning

min/ max  $x_j$  for all  $j$

subject to

$$-b_U \leq Ax \leq -b_L$$

According to the ideas of *Tobias Achterberg*

$$x_L \leq x \leq x_U$$

- Only the first LP subproblem has to be solved from scratch, after that just Phase II is used
- If  $x_j = x_{j,L} / x_{j,U}$  then min / max  $x_j$  can be skipped, resp.
- Sequence of variables: find the non-basic variable that is the closest to its lower / upper bound and has not yet been considered in the pruning step
- LP object is not deallocated between two iteration (memory pool)



# Numerical examples (Separations)



# Liquid-liquid equilibrium (LLE)

Two liquid phases are separated, components are distributed between phases

Solution for the necessary conditions is computed (nonlinear system of equations)



# Liquid-liquid equilibrium (LLE)

$$\sum_i x_i = 1$$

$$\sum_i y_i = 1 \quad i = 1, 2, \dots, C$$

$$\lambda \cdot x_i + (1 - \lambda) \cdot y_i = z_i$$

$$\ln \gamma_i(\mathbf{x}) + \ln x_i = \ln \gamma_i(\mathbf{y}) + \ln y_i$$

$$\sum_{i=1}^{C-1} (x_i - y_i)^2 > \varepsilon$$

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$$\ln \gamma_i(\mathbf{x}) = \frac{\sum_{j=1}^C \tau_{ji} G_{ji} x_j}{\sum_{k=1}^C G_{ki} x_k} + \sum_{j=1}^C \frac{x_j G_{ij}}{\sum_{k=1}^C G_{kj} x_k} \cdot \left( \tau_{ij} - \frac{\sum_{l=1}^C \tau_{lj} G_{lj} x_l}{\sum_{k=1}^C G_{kj} x_k} \right)$$

# LLE results

$$C = 2$$

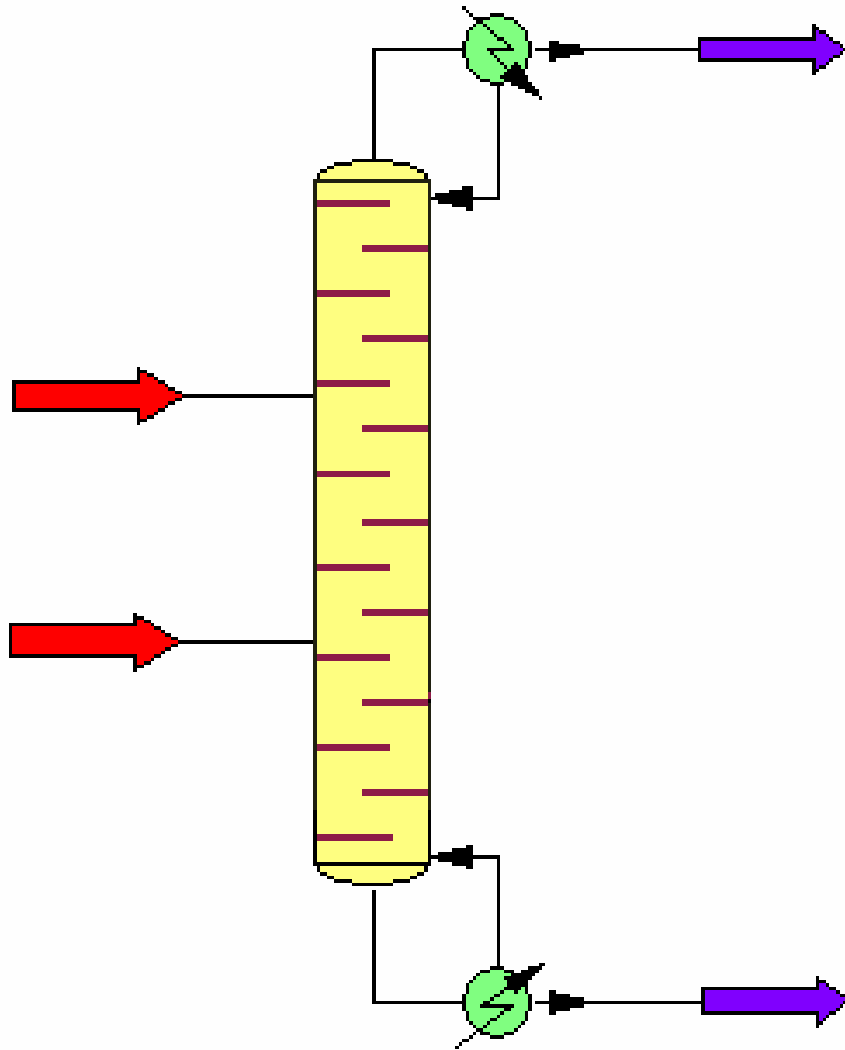
Method	total time [s]	$\frac{\text{total time}}{\text{total time}_{\text{AA/CP}}}$	iterations	$\frac{\text{iterations}}{\text{iterations}_{\text{AA/CP}}}$	cycle time [ $\mu\text{s}$ ]	$\frac{\text{cycle time}}{\text{cycle time}_{\text{AA/CP}}}$
IN/GS	26.1	22.7	120234	85.5	217	0.27
AA/CP	1.15	1.00	1407	1.00	817	1.00

$$C = 3$$

Method	total time [s]	$\frac{\text{total time}}{\text{total time}_{\text{AA/CP}}}$	iterations	$\frac{\text{iterations}}{\text{iterations}_{\text{AA/CP}}}$	cycle time [ms]	$\frac{\text{cycle time}}{\text{cycle time}_{\text{AA/CP}}}$
IN/GS	>318000	$>1.36 \cdot 10^4$	—	—	—	—
AA/CP	23.3	1.00	7715	1.00	3.01	1.00

Baharev'08 Not *state-of-the-art* IN/GB! (C-XSC)

# Distillation columns



# Distillation columns

for  $j = 1, \dots, N$  (*i.e.* for each stage)

$$\sum_i x_{i,j} = 1$$

$$\sum_i y_{i,j} = 1 \quad i = 1, 2, \dots, C$$

$$L_{j-1}x_{i,j-1} + V_{j+1}y_{i,j+1} + F_j z_{i,j} = L_j x_{i,j} + V_j y_{i,j}$$

$$V_j \sum_i \lambda_i y_{i,j} = V_{j+1} \sum_i \lambda_i y_{i,j+1}$$

$$\ln \gamma_{i,j}(\mathbf{x}_j, T_j) + \ln x_{i,j} + (A_i - B_i / (C_i + T_j)) = \ln y_{i,j} + \ln P$$

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$$\ln \gamma_i(\mathbf{x}, T) = -\ln \left( \sum_{a=1}^C x_a \Lambda_{ia} \right) + 1 - \sum_{b=1}^C \frac{x_b \Lambda_{bi}}{\sum_{a=1}^C x_a \Lambda_{ia}} ; \quad \Lambda_{ab}(T) = \frac{V_b^m}{V_a^m} \exp \left( -\frac{k_{ab}}{R_G T} \right)$$

# Results with a single stage

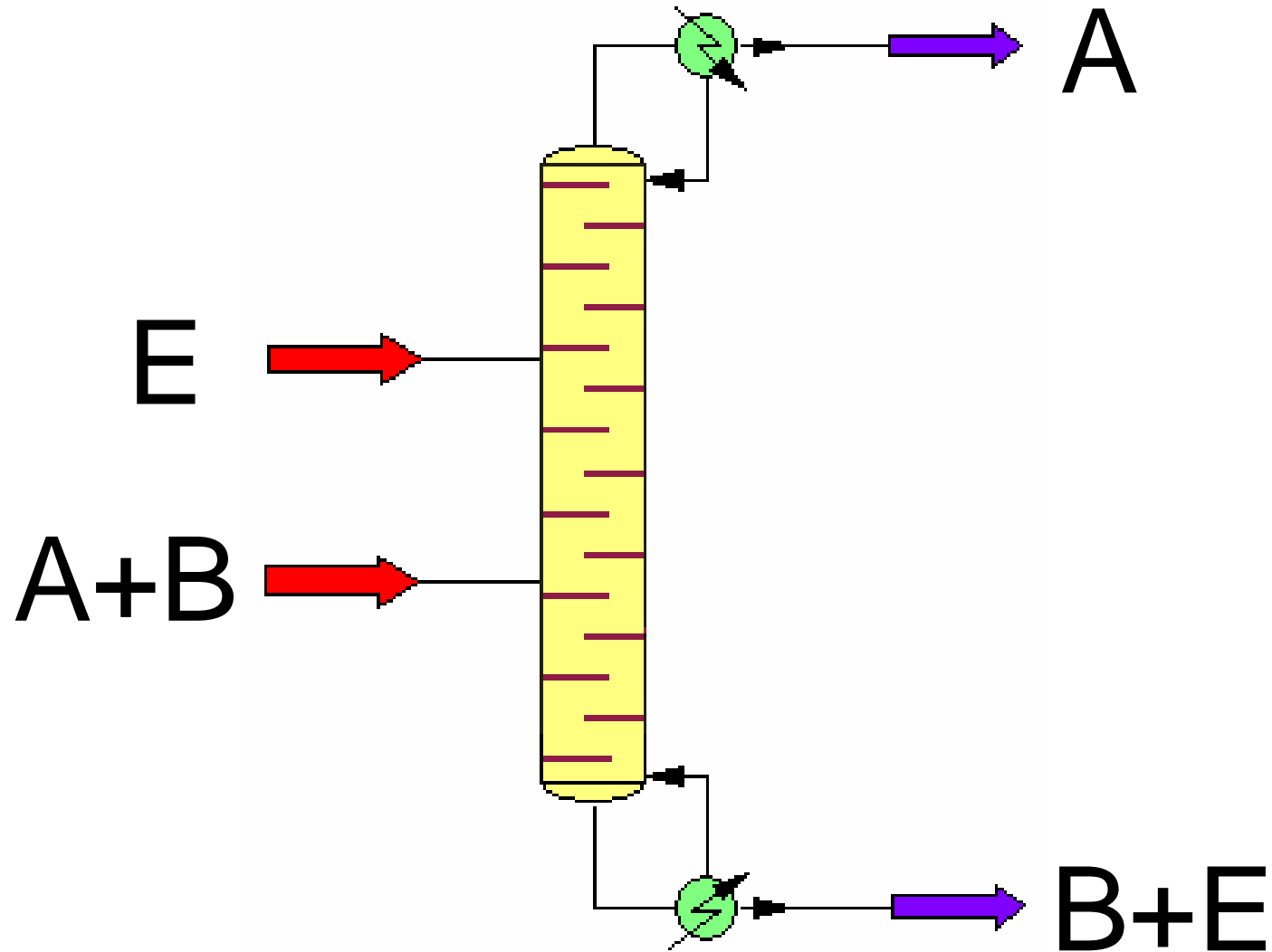
## Linearization techniques ( $N = 1$ )

Method	time [s]	$\frac{\text{total time}_{\text{IN/GS}}}{\text{total time}_{\text{AA/CP}}}$	cycles	$\frac{\text{cycles}_{\text{IN/GS}}}{\text{cycles}_{\text{AA/CP}}}$	cycle time [ms]	$\frac{\text{cycle time}_{\text{IN/GS}}}{\text{cycle time}_{\text{AA/CP}}}$	Cluster boxes
IN/GS	280.7	63.2	271491	160.9	1.03	0.39	3
AA/CP	4.44	1.00	1687	1.00	2.63	1.00	8

## Pruning techniques

Method	time [s]	$\frac{\text{total time}_{\text{IN/GS}}}{\text{total time}_{\text{AA/CP}}}$	cycles	$\frac{\text{cycles}_{\text{IN/GS}}}{\text{cycles}_{\text{AA/CP}}}$	cycle time [ms]	$\frac{\text{cycle time}_{\text{IN/GS}}}{\text{cycle time}_{\text{AA/CP}}}$	Cluster boxes
AA/CP	4.44	4.23	1687	5.68	2.63	0.74	8
AA/LP	1.05	1.00	297	1.00	3.54	1.00	1

# Extractive distillation column



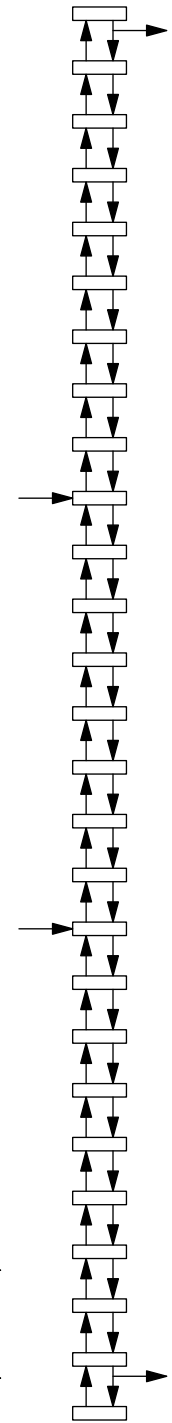
# Distillation column (continued)

Initial intervals are fairly wide

Hundreds of variables

$N$	time (s)	boxes	simplex iterations
12	22.10	19	247522
16	54.15	29	500041
22	92.13	33	709058

Revised implementation: memory pool,  
LP pruning enhanced



# Distillation columns

for  $j = 1, \dots, N$  (i.e. for each stage)

$$\sum_i x_{i,j} = 1$$

$$\sum_i y_{i,j} = 1 \quad i = 1, 2, \dots, C$$

$$L_{j-1}x_{i,j-1} + V_{j+1}y_{i,j+1} + F_j z_{i,j} = L_j x_{i,j} + V_j y_{i,j}$$

$$V_j \sum_i \lambda_i y_{i,j} = V_{j+1} \sum_i \lambda_i y_{i,j+1}$$

$$\ln \gamma_{i,j}(\mathbf{x}_j, T_j) + \ln x_{i,j} + (A_i - B_i / (C_i + T_j)) = \ln y_{i,j} + \ln P$$

---

$$\ln \gamma_i(\mathbf{x}, T) = -\ln \left( \sum_{a=1}^C x_a \Lambda_{ia} \right) + 1 - \sum_{b=1}^C \frac{x_b \Lambda_{bi}}{\sum_{a=1}^C x_a \Lambda_{ia}} ; \quad \Lambda_{ab}(T) = \frac{V_b^m}{V_a^m} \exp \left( -\frac{k_{ab}}{R_G T} \right)$$



# Convex envelopes in LP pruning

$$z = xy$$

$$z \geq x_L y + y_L x - x_L y_L$$

$$z \geq x_U y + y_U x - x_U y_U$$

$$z \leq x_L y + y_U x - x_L y_U$$

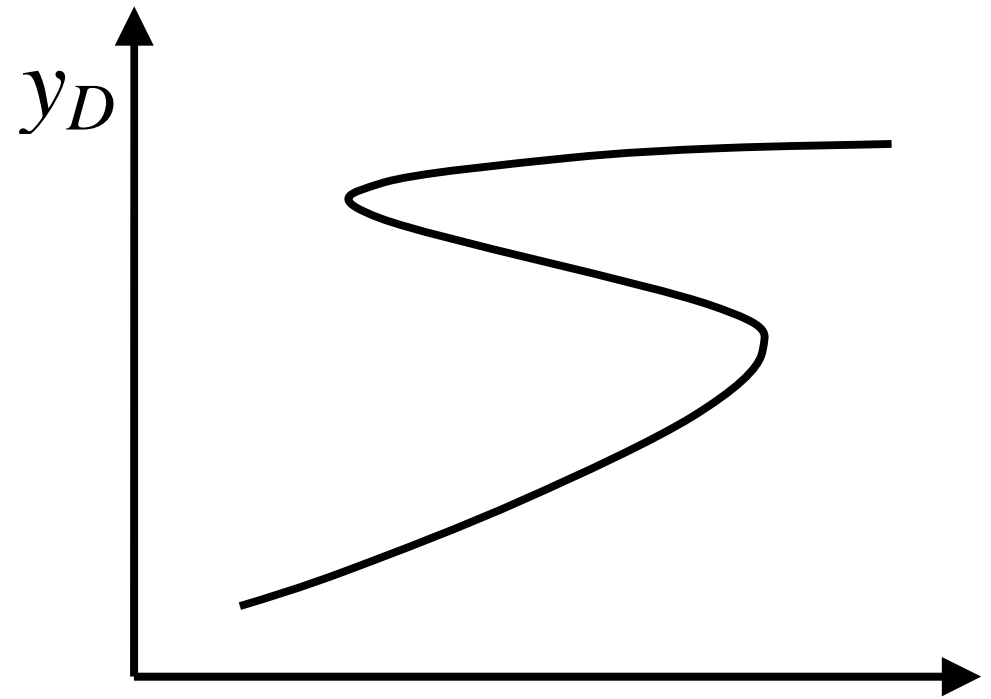
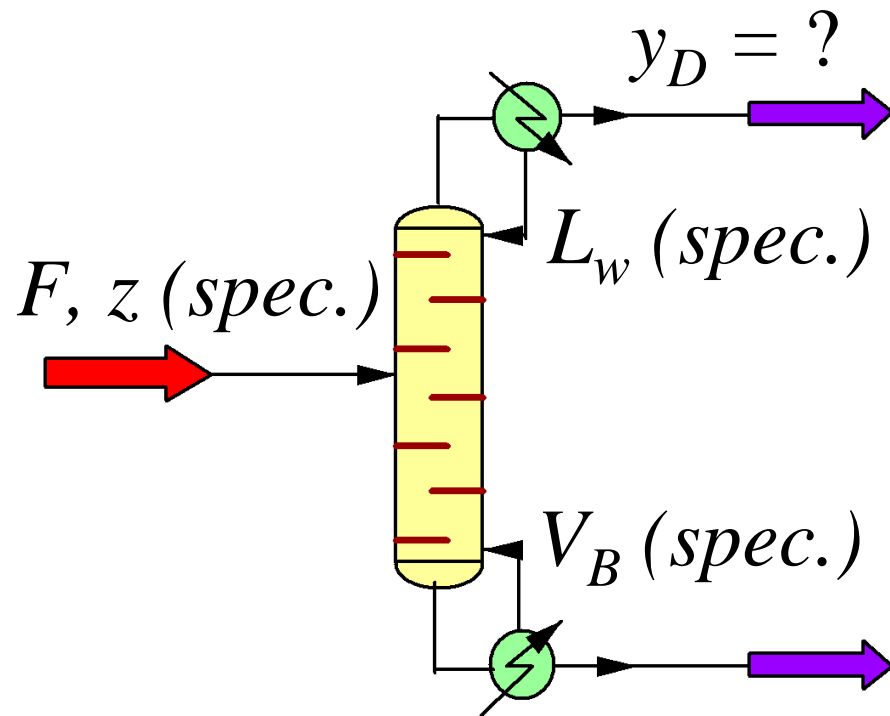
$$z \leq x_U y + y_L x - x_U y_L$$

The initial intervals are varied

	with env.	without env.	with / without env.
time (s)	42.65	9.49	4.50
boxes examined	7	9	0.78
simplex iterations	176411	78803	2.24

	with env.	without env.	with / without env.
time (s)	161.1	54.15	2.97
boxes examined	17	29	0.59
simplex iterations	656571	500041	1.31

# Multiple steady states



Binary mixture ( $C = 2$ )  
Jacobsen'91

Multiple solutions  $L_w$

Instead of the complicated equations of  $\gamma$ :

$$y_j = \alpha x_j / ((\alpha - 1) x_j + 1)$$

# Multiple steady states (results)

Only convex envelopes are used + LP pruning

$L_w$ (kg/min)	$V$ (kmol/min)	number of solutions	time (s)	box	cluster box
58.50	2.0	3	1427	48175	341
96.00	3.0	5	>3600	—	—
98.75	3.0	3	>3600	—	—

## Mixed AA/IA

$L_w$ (kg/min)	$V$ (kmol/min)	number of solutions	time (s)	box	cluster box
58.50	2.0	3	0.22	507	0
96.00	3.0	5	0.25	267	0
98.75	3.0	3	0.20	227	0



# Conclusions



# Conclusions

- Numerical evidence suggests that affine arithmetic is a competing linearization compared to the Interval Newton method
- The revised LP pruning proved to be an efficient pruning technique in certain cases
- The solver written in C++ should be interfaced with a modeling language
- DAG based propagation technique should be studied



# Acknowledgement

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